

Simple I/O

Here we extend our language with a simple I/O mechanism that allows us to initialize variables and examine variables in the executing program:

$$C ::= \text{put}(A) \quad | \quad \text{get}(x)$$

The informal semantics is that 'put' allows us to write an expression to the terminal and 'get' allows us to initialize a declared variable with an integer value read from the terminal.

The formal semantics is as follows:

```
(put(A),State) -->> State :-           %io%  writing
  (A,State) -->> ValA,
  write(A),
  write(' is '),
  writeln(ValA),!.
```

```
(get(X),State) -->> OState :-         %io%  reading
  lookup(X,State,_),
  write('Enter integer value for '),
  write(X),
  write(': '),
  read(Val),
  int(Val),
  bindval(X,Val,State,OState),!.
```

Now we can write programs such as these:

```
?- ['sem-block.pl'].  
% xis.pl compiled 0.01 sec, 7,792 bytes  
% preamble.pl compiled 0.01 sec, 8,956 bytes  
% xis.pl compiled 0.00 sec, 148 bytes  
% sem-block.pl compiled 0.01 sec, 18,284 bytes  
true.
```

```
?- ((var(x) seq get(x) seq put(x)),s) -->> V.  
Enter integer value for x: 3.  
x is 3  
V = push([bind(3, x)], s).
```

```
?- ((var(x) seq get(x) seq put(add(x,1))),s) -->> V.  
Enter integer value for x: 5.  
add(x,1) is 6  
V = push([bind(5, x)], s).
```

```
?-
```

Block Structured Languages

- In most languages a block introduces a new scope allowing for *local variable declarations*. In C blocks are introduced with the curly braces,

```
{
    int x;
}
```

- We can access the values of variables in non-local scope. Consider the following code,

```
{
    int x = 2;
    {
        int y = 3;
        x = y + x; /* accessing the surrounding scope via 'x' */
    }
}
```

- In most languages blocks can be nested \Rightarrow *variable shadowing*.

```
{
    int x = 1;
    {
        int x = 2; /* the original 'x' is no longer visible in this scope */
    }
}
```

Block Structured Languages

Recall that in our simple language we have variable declarations and now we introduce blocks:

$$C ::= \text{var}(x) \mid \text{block}(C)$$

Think of 'block' as 'begin C end' where C could be any command including sequential composition.

Block Structured Languages

Going back to our observations on block structured languages

- Local variable declarations. When we leave a scope with local variables those variables should become undeclared:

```
var(x) seq block( var(y) seq assign(y,1) ) seq assign(y,x)
```

- Non-local side effects. When assigning a value to a variable declared in a surrounding scope we need to update the value of that variable, the value printed out for x should be 2:

```
var(x) seq assign(x,1) seq block( assign(x,2) ) seq put(x)
```

- Variable shadowing. Redeclaring a variable in a nested scope with the same name as a variable in the outer scope makes the variable in the outer scope unavailable, the value printed out for x should be 1:

```
var(x) seq assign(x,1) seq block( var(x) seq assign(x,2) ) seq put(x)
```

Block Structured Languages

Formal Semantics:

```
(var(X),State) -->> OState :-           % decl,  
    declarevar(X,State,OState),!.  
  
(assign(X,A),State) -->> OState :-     % assignment  
    lookup(X,State,_),                 % only allowed to assign to variables  
    (A,State) -->> ValA,                % that have been declared  
    bindval(X,ValA,State,OState),!.  
  
(block(C),State) -->> OState :-       %block%    block statement  
    pushenv(State,LocalState),  
    (C,LocalState) -->> S,  
    popenv(S,OState),!.
```

Note: Each block now pushes its own binding environment on an environment stack.

Note: The new semantic procedures 'declarevar' and 'bindval' are necessary because declaring and binding is done differently with nested scopes.

Block Structured Languages

Semantic procedures 'pushenv' and 'popenv':

```
%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%  
% the predicate 'pushenv(+State,-FinalState)' pushes  
% a new binding term list on the stack  
:- dynamic pushenv/2.
```

```
pushenv(S,env([],S)) :- !.
```

```
%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%  
% the predicate 'popenv(+State,-FinalState)' pops  
% a binding term list off the stack  
:- dynamic popenv/2.
```

```
popenv(env(_,S),S) :- !.
```


Block Structured Languages

Looking up variable bindings in a stack of binding environments:

```
% the predicate 'lookup(+Variable,+State,-Value)' looks up
% the variable in the state and returns its bound value.
:- dynamic lookup/3.                % modifiable predicate

lookup(_,s0,_) :- fail.

lookup(X,env([],S),Val) :-
    lookup(X,S,Val),!.

lookup(X,env([bind(Val,X)|_],_),Val).

lookup(X,env([_|Rest],S),Val) :-
    lookup(X,env(Rest,S),Val),!.
```

Block Structured Languages

Semantic procedure 'declarevar':

```
% the predicate 'declarevar(+Variable,+State,-FinalState)' declares
% a variable by inserting a new binding term into the current
% environment.
:- dynamic declarevar/3.                                % modifiable predicate

declarevar(X,S,env([bind(0,X)],S)) :-
    atom(S),!.

declarevar(X,env(L,S),env([bind(0,X)|L],S)) :- !.
```

Block Structured Languages

Semantic procedure 'bindval':

```
% the predicate 'bindval(+Variable,+Value,+State,-FinalState)' updates
% a binding term in the state.  this update is done "in place"
% in order to support global variables.  the predicate has to
% search both the binding list and the stack of binding
% lists.
:- dynamic bindval/4.                                % modifiable predicate

bindval(_,_ ,s0,_ ) :-
    fail.

bindval(X,Val,env([],S),env([],NewS)) :-
    bindval(X,Val,S,NewS),!.

bindval(X,Val,env([bind(_,X)|L],S),env([bind(Val,X)|L],S)),!.

bindval(X,Val,env([bind(V,Y)|L],S),env([bind(V,Y)|NewL],NewS)) :-
    bindval(X,Val,env(L,S),env(NewL,NewS)),!.
```

Block Structured Languages

```
?- ['sem-block.pl'].
%   xis.pl compiled 0.00 sec, 7,792 bytes
%  preamble.pl compiled 0.00 sec, 8,956 bytes
%   xis.pl compiled 0.00 sec, 148 bytes
% sem-block.pl compiled 0.00 sec, 18,192 bytes
true.

?- ((var(x) seq block( var(y) seq assign(y,1) ) seq assign(y,x)),s) --> V.
false.

?- ((var(x) seq assign(x,1) seq block( assign(x,2) ) seq put(x)),s) --> V.
x is 2
V = env([bind(2, x)], s).

?- (( var(x) seq assign(x,1) seq block( var(x) seq assign(x,2) ) seq put(x)),s) --> V.
x is 1
V = env([bind(1, x)], s).

?-
```