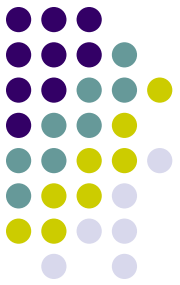




Scope & Symbol Table

- Most modern programming languages have some notion of scope.
- Scope defines the “lifetime” of a program symbol.
- If a symbol is no longer accessible then we say that it is “out of scope.”
- The simplest scope is the “block scope.”
- With scope we need a notion of variable declaration which allows us to assert in which scope the variable is visible or accessible.

Read Chap 7



Cuppa2

- We extend our Cuppa1 language with variable declarations of the form

declare $x = 10$;

- Declares the variable x in the current scope and initializes it to the value 10
- If the current scope is the global (outermost) scope then we call x a “global” variable.

Cuppa2 Grammar



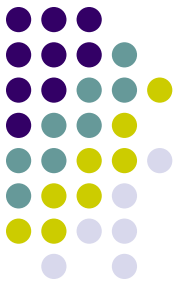
Listing 7.1: Grammar for the Cuppa2 language.

```
1 stmt_list : (stmt)*
2
3 stmt : declare ID (= exp)? ;?
4       | ID = exp ;?
5       | get ID ;?
6       | put exp ;?
7       | while \( exp \) stmt
8       | if \( exp \) stmt (else stmt)?
9       | \{ stmt_list \}
10
11 exp : exp_low
12 exp_low : exp_med ((= | =<) exp_med)*
13 exp_med : exp_high ((+ | -) exp_high)*
14 exp_high : primary ((\* | /) primary)*
15
16 primary : INTEGER
17          | ID
18          | \( exp \)
19          | - primary
20          | not primary
21
22 ID : <any valid variable name>
23 INTEGER : <any valid integer number>
```

```
stmt : declare ID (= exp)? ;?
```

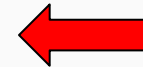
Notice that the initializer for the declaration is optional.

Cuppa2 Frontend

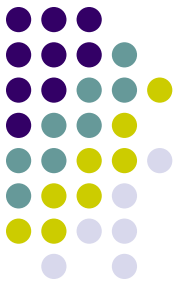


cuppa2_fe.py

```
def stmt(stream):
    token = stream.pointer()
    if token.type in ['DECLARE']:
        stream.match('DECLARE')
        id_tk = stream.match('ID')
        if stream.pointer().type in ['ASSIGN']:
            stream.match('ASSIGN')
            e = exp(stream)
        else:
            # if no initializer assume default value
            e = ('INTEGER', 0)
        if stream.pointer().type in ['SEMI']:
            stream.match('SEMI')
        return ('DECLARE', ('ID', id_tk.value), e)
    elif token.type in ['ID']: ...
```



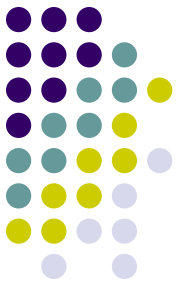
The relevant piece of code in the frontend.



Cuppa2

- We can now write properly scoped programs
- Consider:

```
declare x = 1;
{
  declare x = 2;
  put x;
}
{
  declare x = 3;
  put x;
}
put x;
```

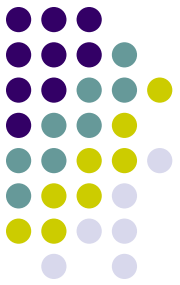


Variable Shadowing

- An issue with scoped declarations is that inner declarations can “overshadow” outer declarations
- Consider:

```
declare x = 2;
{
  declare x = 3;
  {
    declare y = x + 2;
    put y;
  }
}
```

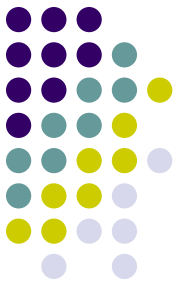
What is the output of the program once it is run?



Variable update

- A variable update can be outside of our current scope.
- Consider

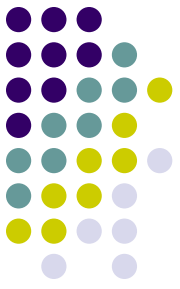
```
declare x = 2;
{
  declare y = 3;
  x = y + x;
  put x;
}
put x;
```



Symbol Tables

- To deal with programs like that we need something more sophisticated for variable lookup than a dictionary.
 - ☞ *a dictionary stack*
- This stack needs to be able to support the following functionality
 - Declare a variable (insertion)
 - Lookup a variable
 - Update a variable value

Semantic Rules for Variable Declarations



- Here are the rules which we informally used in the previous examples:
 - The ‘declare’ statement inserts a variable declaration into the current scope
 - a variable lookup returns a variable value from the current scope or the surrounding scopes
 - Every variable needs to be declared before use
 - No variable can be declared more than once in the current scope.

Symbol Tables

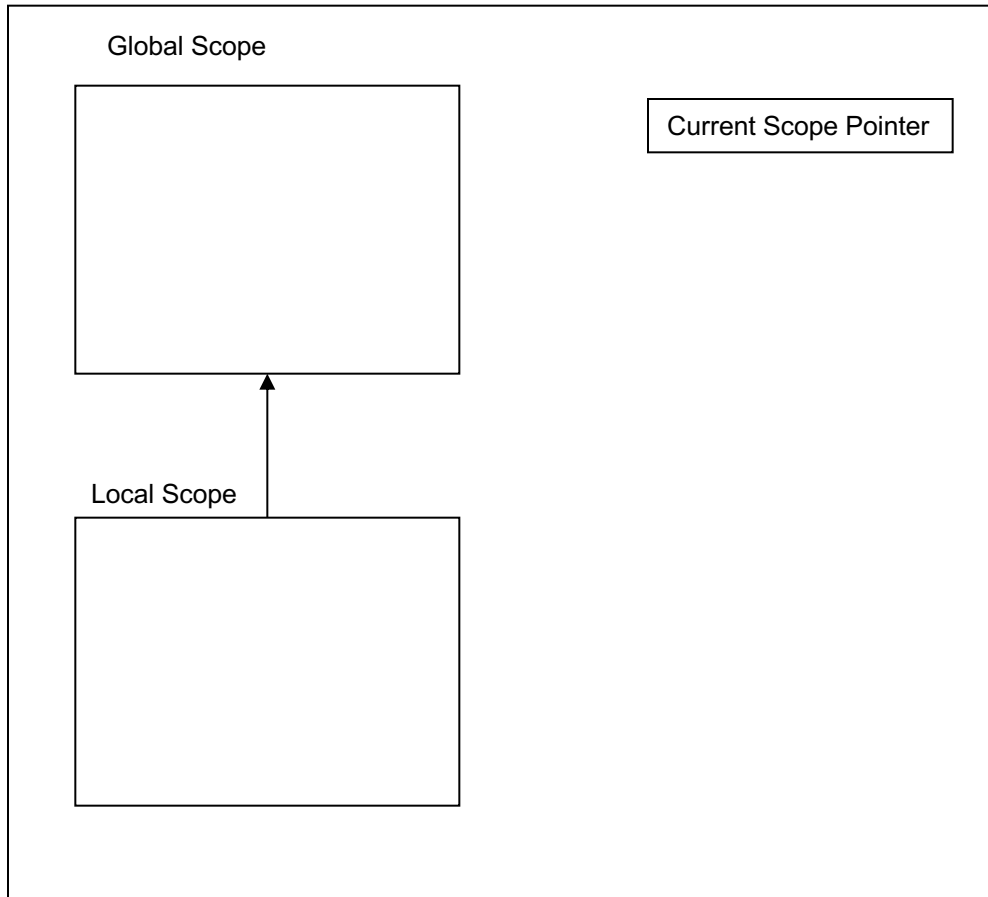


- Design:
 - we have a class **SymTab** that:
 - Holds a stack of scopes
 - `scoped_symtab`
 - Defines the interface to the symbol table
 - `push_scope`, `pop_scope`, `declare_sym`, *etc*
 - By default, SymTab is initialized with a single scope on the stack – *the global scope*.

Symbol Tables

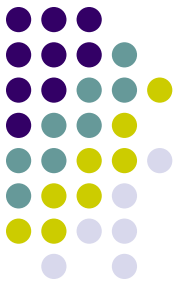


Symbol Table

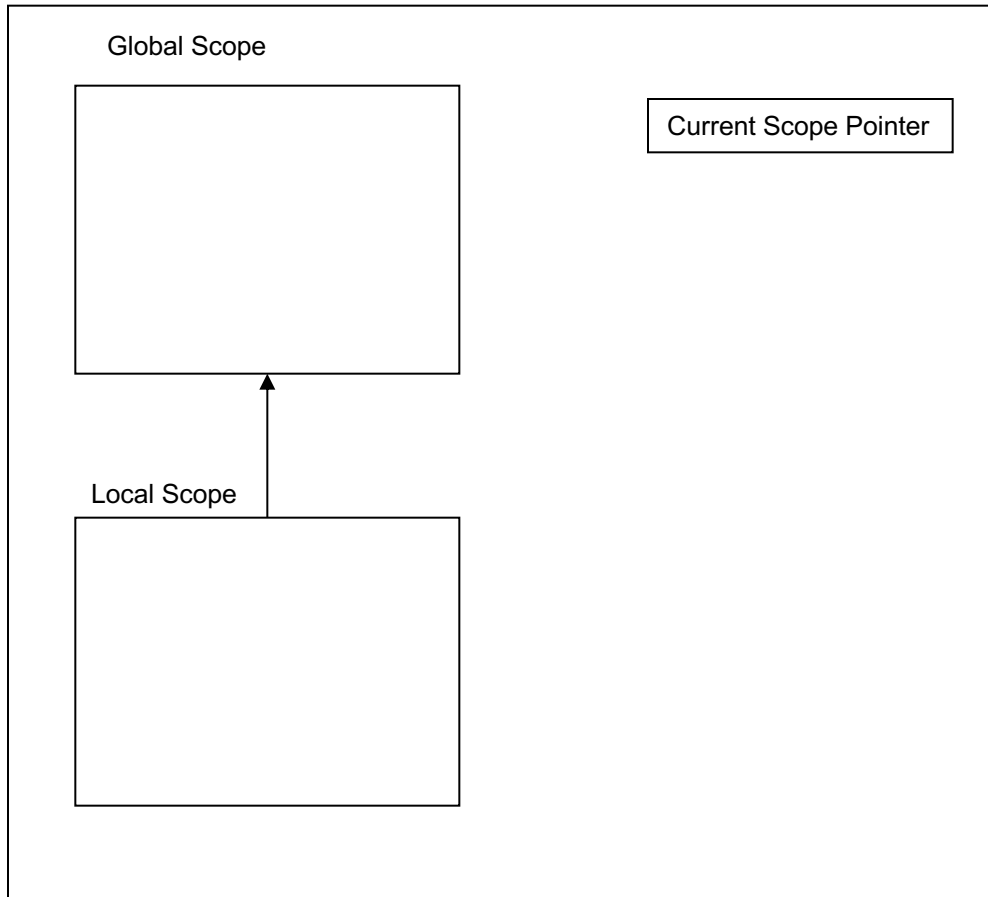


```
declare x = 2;  
{  
    declare y = 3;  
    x = y + x;  
    put x;  
}  
put x;
```

Symbol Tables

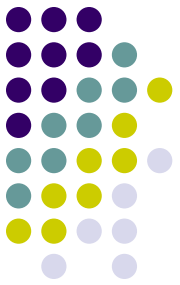


Symbol Table

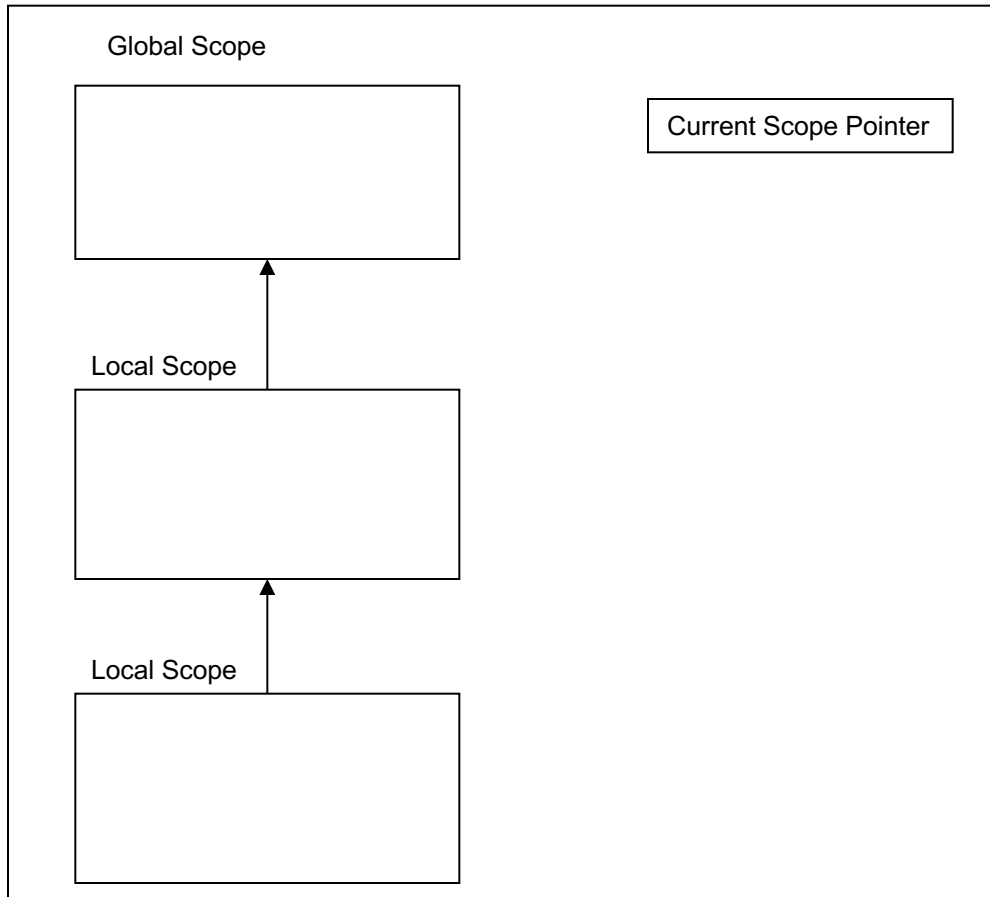


```
declare x;  
get x;  
If (0 <= x)  
{  
    declare i = x;  
    put i;  
}  
else  
{  
    declare j = -1 * x;  
    put j;  
}  
put x;
```

Symbol Tables



Symbol Table



```
declare x = 2;
{
  declare x = 3;
  {
    declare y = x + 2;
    put y;
  }
}
```

Symbol

cuppa2_syntab.py

```
CURR_SCOPE = 0

class SymTab:

    #-----
    def __init__(self):
        # global scope dictionary must always be present
        self.scoped_syntab = [{}]

    #-----
    def push_scope(self):
        # push a new dictionary onto the stack - stack grows to the left
        self.scoped_syntab.insert(CURR_SCOPE, {})

    #-----
    def pop_scope(self):
        # pop the left most dictionary off the stack
        if len(self.scoped_syntab) == 1:
            raise ValueError("cannot pop the global scope")
        else:
            self.scoped_syntab.pop(CURR_SCOPE)

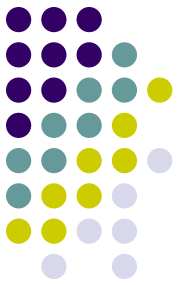
    #-----
    def declare_sym(self, sym, init):
        # declare the symbol in the current scope: dict @ position 0
        ...

    #-----
    def lookup_sym(self, sym):
        # find the first occurrence of sym in the syntab stack
        # and return the associated value
        ...

    #-----
    def update_sym(self, sym, val):
        # find the first occurrence of sym in the syntab stack
        # and update the associated value
        ...
```



Symbol Tables



```
def declare_sym(self, sym, init):
    # declare the symbol in the current scope: dict @ position 0

    # first we need to check whether the symbol was already declared
    # at this scope
    if sym in self.scoped_syntab[Curr_Scope]:
        raise ValueError("symbol {} already declared".format(sym))

    # enter the symbol in the current scope
    scope_dict = self.scoped_syntab[Curr_Scope]
    scope_dict[sym] = init
```

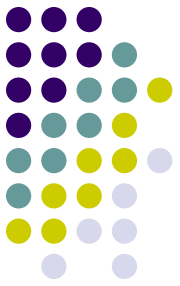
```
def lookup_sym(self, sym):
    # find the first occurrence of sym in the syntab stack
    # and return the associated value

    n_scopes = len(self.scoped_syntab)

    for scope in range(n_scopes):
        if sym in self.scoped_syntab[scope]:
            val = self.scoped_syntab[scope].get(sym)
            return val

    # not found
    raise ValueError("{} was not declared".format(sym))
```

Symbol Tables



```
def update_sym(self, sym, val):
    # find the first occurrence of sym in the symtab stack
    # and update the associated value

    n_scopes = len(self.scoped_symtab)

    for scope in range(n_scopes):
        if sym in self.scoped_symtab[scope]:
            scope_dict = self.scoped_symtab[scope]
            scope_dict[sym] = val
            return

    # not found
    raise ValueError("{} was not declared".format(sym))
```

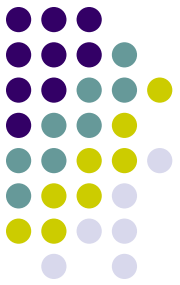

Interpret Walker

Note: Same as Cuppa1 interpreter except for the addition of the declaration statement and additional functionality in block statements and variable expressions.

cuppa2_interp_walk.py


```
def walk(node):
    # node format: (TYPE, [child1[, child2[, ...]])
    type = node[0]
    if type in dispatch:
        node_function = dispatch[type]
        return node_function(node)
    else:
        raise ValueError("walk: unknown tree node type: " + type)



# a dictionary to associate tree nodes with node functions
dispatch = {
    'STMTLIST' : stmtlist,
    'DECLARE'  : declare_stmt,
    'ASSIGN'   : assign_stmt,
    'GET'      : get_stmt,
    'PUT'      : put_stmt,
    'WHILE'    : while_stmt,
    'IF'       : if_stmt,
    'NIL'      : nil,
    'BLOCK'    : block_stmt,
    'INTEGER'  : integer_exp,
    'ID'       : id_exp,
    'PAREN'    : paren_exp,
    'PLUS'     : plus_exp,
    'MINUS'    : minus_exp,
    'MUL'      : mul_exp,
    'DIV'      : div_exp,
    'EQ'       : eq_exp,
    'LE'       : le_exp,
    'UMINUS'   : uminus_exp,
    'NOT'      : not_exp
}
```




Interpret Walker

```
def declare_stmt(node):  
  
    (DECLARE, (ID, name), exp) = node  
  
    value = walk(exp)  
    symbol_table.declare_sym(name, value)  
  
    return None
```




```
def block_stmt(node):  
  
    (BLOCK, stmt_list) = node  
  
    symbol_table.push_scope()   
    walk(stmt_list)  
    symbol_table.pop_scope()   
  
    return None
```

```
def assign_stmt(node):  
  
    (ASSIGN, (ID, name), exp) = node  
  
    value = walk(exp)  
    symbol_table.update_sym(name, value)  
  
    return None
```

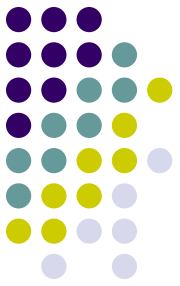


```
def get_stmt(node):  
  
    (GET, (ID, name)) = node  
  
    s = input("Value for " + name + '? ')  
  
    try:  
        value = int(s)  
    except ValueError:  
        raise ValueError("expected an integer value for " + name)  
  
    symbol_table.update_sym(name, value)  
  
    return None
```

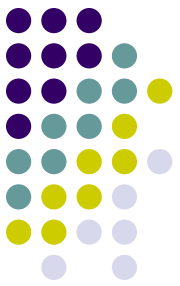


That's it – everything else is the same as the Cuppa1 interpreter!

Syntactic vs Semantic Errors



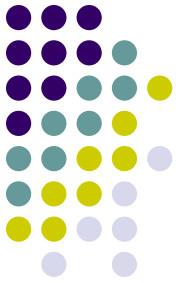
- Grammars allow us to construct parsers that recognize the syntactic structure of languages.
- Any program that does not conform to the structure prescribed by the grammar is rejected by the parser.
- We call those errors “syntactic errors.”



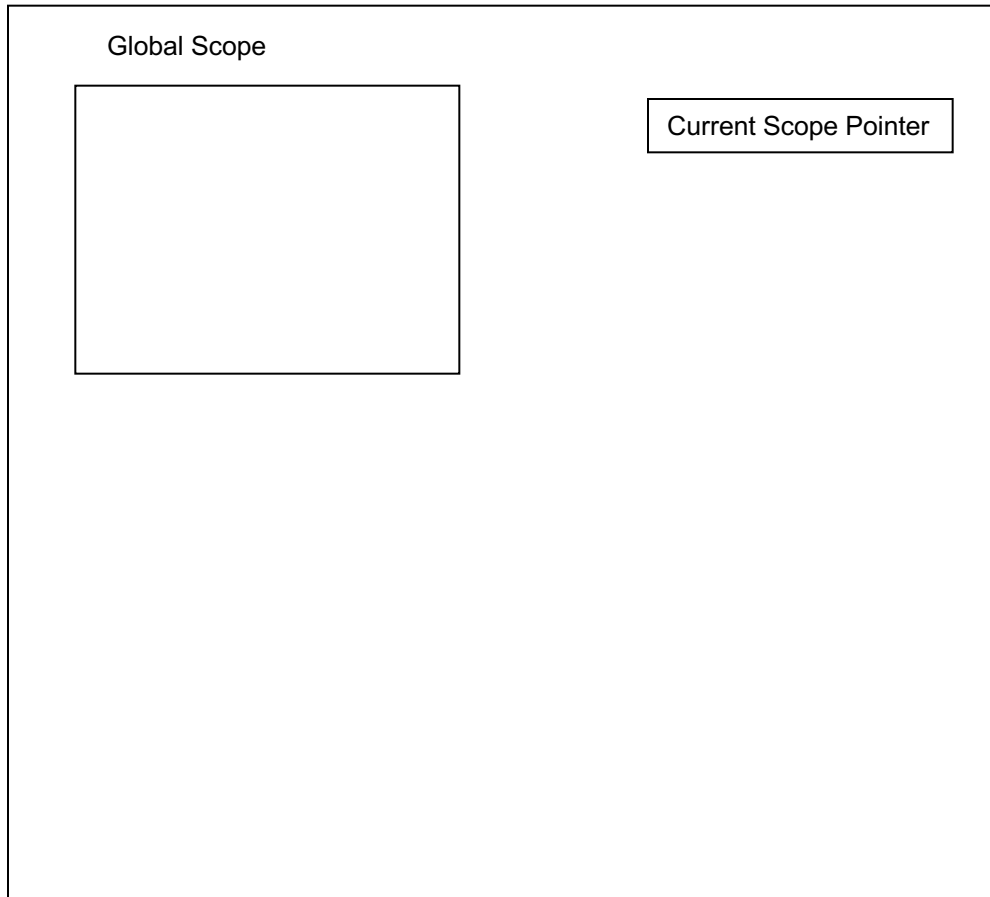
Syntactic vs Semantic Errors

- Semantic errors are errors in the behavior of the program and cannot be detected by the parser.
- **Programs with semantic errors are usually syntactically correct**
- A certain class of these semantic errors can be caught by the interpreter/compiler. Consider:
 declare x = 10;
 put x + 1;
 declare x = 20;
 put x + 2;
- Here we are redeclaring the variable ‘x’ which is not legal in many programming languages.
- Many other semantic errors cannot be detected by the interpreter/compiler and show up as “bugs” in the program.

Symbol Tables

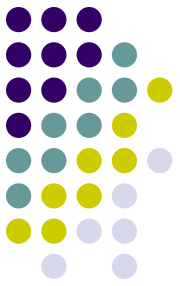


Symbol Table

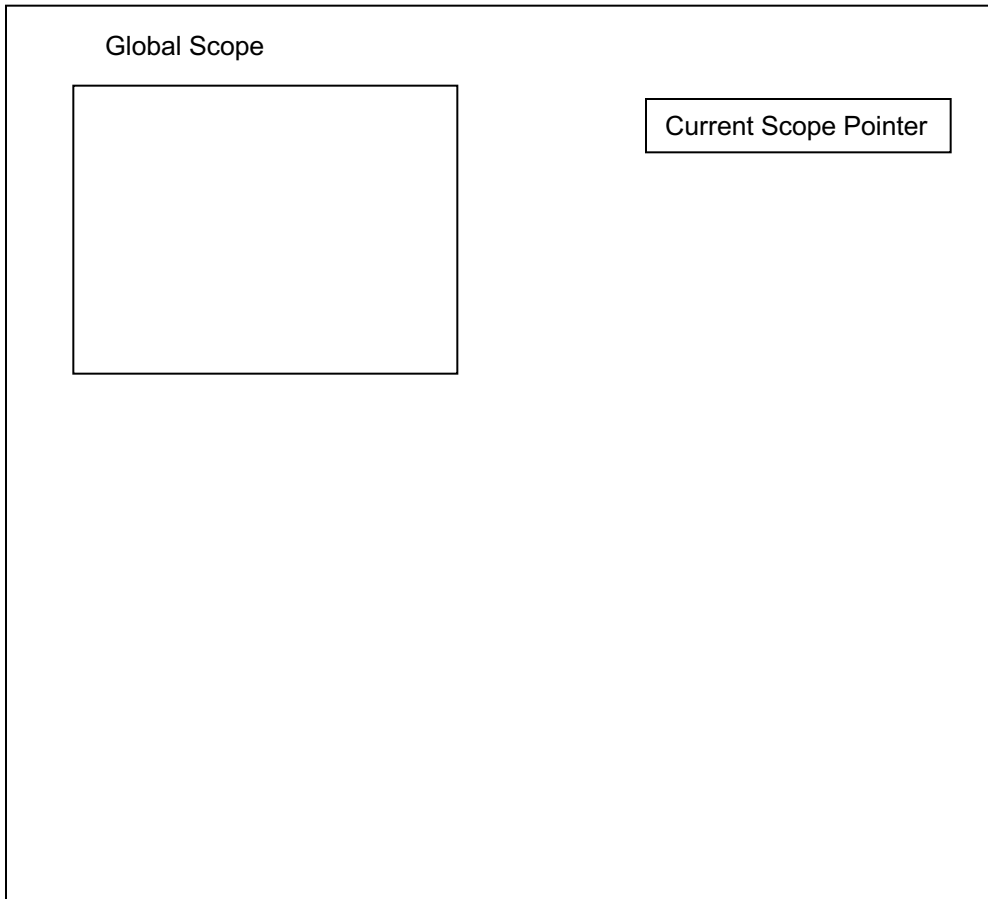


```
declare x = 10;  
put x + 1;  
declare x = 20;  
put x + 2;
```

Symbol Tables



Symbol Table



```
x = x + 1;  
put x;
```