

# Defining Language TWO

- Extend Language ONE with:
  - Variables
  - `let` expression for assigning values to them

# Reading

- Section 23.3

# TWO: Syntax

TWO:

$\langle exp \rangle^* ::= \langle exp \rangle + \langle mulexp \rangle \mid \langle mulexp \rangle$   
 $\langle mulexp \rangle ::= \langle mulexp \rangle * \langle rootexp \rangle \mid \langle rootexp \rangle$   
 $\langle rootexp \rangle ::= \mathbf{let\ val} \langle variable \rangle = \langle exp \rangle \mathbf{in} \langle exp \rangle \mathbf{end}$   
 $\quad \mid (\langle exp \rangle) \mid \langle variable \rangle \mid \langle constant \rangle$

- A sample Language TWO expression:  
`let val y = 3 in y*y end`
- What does the parse tree for the above expression look like?
- Notice that in TWO assignments are expressions and the variables bindings are only valid in the scope of the second expression of the let expression.

# TWO: Abstract Syntax

Additional abstract syntax nodes for language TWO:

- (1) var(X) dereferences a variable X
- (2) let(X,E1,E2) binds the variable X to expression E1 in the context of expression E2.

Example: the TWO program

```
let val y = 3 in y*y end
```

will result in the AST

```
let (y, const(3), times(var(y), var(y)))
```

# From Parse Tree to Prolog AST

- Consider: 2 \* let x = 5 in 1+x end
  - Parse tree?
  - AST?
  - Prolog AST?

# TWO: Semantics

In order to provide semantics we need to remember the values assigned to variables -- binding environments (fancy word for dictionary!)

In our case, for the Prolog based semantics, we let the terms `bind(X,K)` represent the binding of variable `X` to value `K`. A context is simply a list of these binding terms:

```
[bind(y, 3), bind(q, 20), bind(z, 5)]
```

Given this binding structure, we can write a predicate, `lookup/3`, that returns a variable binding for a particular `Var`

```
lookup(Var, [bind(Var, Value) | _ ], Value) .  
lookup(Var, [ _ | Rest], Value) :- lookup(Var, Rest, Value) .
```

Finds the most recent binding of variable `Var` if there is one.

# TWO: Prolog Interpreter

```
val2 (plus (X,Y) ,B,Value) :-  
    val2 (X,B,XValue) ,  
    val2 (Y,B,YValue) ,  
    Value is XValue + YValue.  
  
val2 (times (X,Y) ,B,Value) :-  
    val2 (X,B,XValue) ,  
    val2 (Y,B,YValue) ,  
    Value is XValue * YValue.  
  
val2 (const (X) ,_,X) .  
  
val2 (var (X) ,B,Value) :-  
    lookup (X,B,Value) .  
  
val2 (let (X,Exp1,Exp2) ,B,Value) :-  
    val2 (Exp1,B,XValue) ,  
    val2 (Exp2, [bind (X,XValue) | B] ,Value) .
```

val2 / 3 - interpretation predicate  
first argument: AST  
second argument: binding env;  
third argument: semantic value.

# Examples

```
let val y = 3 in y*y end
```

```
?- val2(let(y,const(3),times(var(y),var(y))),[ ],X).
```

```
X = 9
```

```
Yes
```

```
let val y = 3 in
  let val x = y*y in
    x*x
  end
end
```

```
let val y = 1 in
  let val y = 2 in
    y
  end
end
```



# Exercises

- Use the semantics of TWO to show the following:
  - Assume that the context  $B = [\text{bind}(y,3)]$  then the semantic value of  $'2*y'$  is 6
  - The semantic value of  $'2 * \text{let } x = 3 \text{ in } x * x \text{ end}'$  is 18
  - The semantic value of  $'\text{let } x = 1 \text{ in let } y = x + 1 \text{ in } y \text{ end end}'$  is 2

# Exercises

- Use the semantics to compute the meaning of the following expressions in TWO (use the rules given in the notes, the book has many typos):

1) `let val y = 3 in 2*y end`

2) `let val y = 1 in  
 let val y = 2 in  
 y  
 end  
end`

Note: first construct an abstract syntax tree, then give the representation in Prolog notation, and then show the computation in our semantics.