Modern programming languages have many different classes of Variables, e.g.

(1) Global variables
(2) Parameters
(3) (function) local variables (also called automatic or activation-specific)
(4) (object-oriented) member variables
(5) Etc.

It is the job of the language system to keep track of the values of these variables during the runtime of a program.

⇒ The language system accomplishes this by binding a variable to a memory location and then uses that memory location to store the value of the variable.
Reading

- Chap 12 in MPL
Memory Locations for Variables

In imperative programs this is a fairly transparent process - the assignment operator mimics what happens at the hardware level - namely, the updating of memory cells.

In functional and logic programming languages this is often not so obvious, since there is no global state, but still, variables are bound to memory locations.
In order to track variables for functions, compilers use a data structure called activation record - collects all the variables belonging to one function into this structure.

Example: FORTRAN

FUNCTION AVG (M,N)
SUM = M + N
AVG = SUM/2.0
RETURN
END

Note: Activation Records are often called Frames
Note: Non-recursive languages such as FORTRAN keep a single activation record per function in the program.

Recursive languages (ML, Java, C, C++, etc) keep a stack of activation records; one per function call.
The Runtime Stack

Code:
(Main)
...
(function code)
...
RETURN

Global Data:
(Stack pointer)

Activation Record
<local vars>
<return addr>
<next record>

Activation Record
<local vars>
<return addr>
<next record>

Runtime Stack
We are evaluating \textbf{fact(3)}. This shows the contents of memory just before the recursive call that creates a second activation.

```java
int fact(int n) {
    int result;
    if (n<2) result = 1;
    else result = n * fact(n-1);
    return result;
}
```
This shows the contents of memory just before the third activation.

```java
int fact(int n) {
    int result;
    if (n<2) result = 1;
    else result = n * fact(n-1);
    return result;
}
```
int fact(int n) {
    int result;
    if (n<2) result = 1;
    else result = n * fact(n-1);
    return result;
}

This shows the contents of memory just before the third activation returns.
The second activation is about to return.

```java
int fact(int n) {
    int result;
    if (n<2) result = 1;
    else result = n * fact(n-1);
    return result;
}
```
The first activation is about to return with the result \textbf{fact(3) = 6}.

```java
int fact(int n) {
    int result;
    if (n<2) result = 1;
    else result = n * fact(n-1);
    return result;
}
```